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Application transmittal

9/19/00

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# UTILITY PATENT APPLICATION TRANSMITTAL

(only for new non-provisional applications under 37 CFR 1.53(b))

Attorney Docket No. **Brown 2000-0315**

First Named Inventor or Application Identifier **John C. Brown**

Title **Controlled Transmissions Across Packet Networks**

Express Mail Label no. **EJ936900221US**

To: **Assistant Commissioner for Patents  
Box Patent Application  
Washington D.C. 20231**

## APPLICATION ELEMENTS

## ACCOMPANYING APPLICATION PARTS

- ☒ Fee Transmittal Form (original and duplicate)
- ☒ Specification **Total Pages 18**  
 title  
 cross reference to related applications (e.g. provisional application)  
 background  
 summary  
 brief description of the drawings (if filed)  
 detailed description  
 claims  
 abstract
- ☒ Drawing(s) **Total Pages 5**
- ☒ Declaration **Total Pages 2**  
 a. ☐ Newly executed  
 b. ☐ Copy from a prior application (37 CFR 1.63(d))  
 (for continuations/divisionals with section below filled out)  
☐ Deletion of Inventor(s) Signed Statement attached deleting  
 inventor(s) named in the prior application. 37 CFR 163 (d)(2)  
 and 1.33(b).  
☐ Incorporation by reference (usable if Declaration is a copy):  
 The entire disclosure of the prior application, from which a copy of the oath or declaration  
 is supplied, is considered as being part of the disclosure of the accompanying application  
 is hereby incorporated by reference herein.  
☐ Other

- ☒ Assignment
- ☒ Recordation form
- ☒ Power of Attorney
- ☒ Postcard
- ☐ Small entity statement
- ☐ Certified copy of priority documents
- ☐ Information disclosure statement
- ☐ Copies of IDS citations
- ☐ 37 CFR 3.73(b) Statement
- ☒ check
- ☐ Other

If a CONTINUING APPLICATION, check appropriate box and supply the requisite information:

- ☐ Continuation ☐ Divisional ☐ Continuation-in-part (CIP) of prior Application No:

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☐ Customer Number or Bar Code Label

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Reg. No. **26,844**

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Signature

*Henry Brendzel*

Date

9/18/00

I hereby certify that this Application is being deposited with the United States Postal Service "Express Mail Post Office to Addressee" service under 37 CFR 1.10 on the date indicated above and is addressed to the Assistant Commissioner for Patents, Washington D.C. 20231.


9/18/00  
Date of Deposit

**Henry Brendzel**  
(Printed Name of Person Mailing Paper)

*Henry Brendzel*  
(Signature of Person Mailing Paper)

<p align="center"><b>FEE TRANSMITTAL</b></p> <p><i>Patent Fees are subject to annual revisions on October 1.          These are the fees effective November 10, 1998          Small entity payments must be supported by a small entity statement,          Otherwise, large entity fees must be paid. See Forms PTO/SB/09-12.</i></p>		PTO/SB/17 Modified 12/98 - Henry Brenzetti	
		Complete if Known	
		Application Number	
		Filing Date	
		First Named Inventor	John C. Brown
		Examiner Name	
		Group/Art Unit	
TOTAL AMOUNT OF PAYMENT	(\$) <b> 748</b>	Attorney Docket ID	Brown 2000-0315

METHOD OF PAYMENT (check one)						FEE CALCULATION (continued)	
1. The Commissioner is hereby authorized to charge indicated fees and other underpayments, and credit overpayments to:						3. ADDITIONAL FEES	
Deposit Account Number  <b>500732</b>		Deposit Account Name  <b>Henry T. Brendzel</b>					
<input checked="" type="checkbox"/> Charge any additional Fee Required under 37 CFR 1.16 and 1.17		<input type="checkbox"/> Charge the Issue Fee Set in 37 CFR 1.18 at the Mailing of the Notice of Allowance					
2. <input type="checkbox"/> Payment enclosed:							
<input type="checkbox"/> Check		<input type="checkbox"/> Money Order		<input type="checkbox"/> Other			
FEE CALCULATION							
1. FILING FEE							
Fee Description					Fee Paid		
Utility Filing Fee .....					690		
Design Filing Fee .....							
Plant Filing Fee .....							
Reissue Filing Fee .....							
Provisional Filing Fee .....							
SUBTOTAL (1) (\$)					690		
2. CLAIMS							
	Claims remaining	Highest Paid	Extra	Rate	Amount		
Total Claims	21	20	1	18	18		
Independent Claims	2	3	0	78	0		
Multiply Dependent Claims	<input type="checkbox"/>			260	0		
SUBTOTAL (2) (\$)					18		
						Other fee Other fee	
SUBTOTAL (2) (\$)						40	
SUBTOTAL (2) (\$)						40	

<b>SUBMITTED BY</b>				<b>Complete (if applicable)</b>	
Typed of Printed Name	<b>Henry T. Brendzel</b>			Reg. Number	<b>26,844</b>
Signature		Date	<b>7/18/00</b>	Deposit Account User ID	

[illegible]

When a call from user 11 to user 21 is assigned to traverse packet network 300 that employs, for example, the IP protocol, central office 12 sends signaling information to VoP gateway 10 that couples network 100 to packet network 300. Gateway 10 translates and converts the received signaling information to a chosen signaling format, for example Media Gateway Control Protocol (MGCP) over IP, and forwards the signaling packets to call agent 15. The signaling packets contain information such as the identity of the called party and the identity of the calling party. Call agent 15 queries database 16 (with the destination of called party 21) to identify an appropriate call agent for completing the connection, and receives the IP address of PacTel call agent 25. Call agent 15 then sends an Initial Address Message (IAM) to call agent 25, requesting the IP address of the appropriate VoP gateway for completing the call. Call agent 25 queries its database (26), obtains the IP address of VoP 20, and forwards that information in an Address Complete Message (ACM) to call agent 15. The communication path

between the call agents is not shown, for sake of clarity. The communication itself can employ the Bearer Independent Call Control (BICC) protocol. The IP address of VoP gateway 10 is communicated to VoP gateway 20 by call agent 25, the IP address of VoP gateway 20 is communicated to VoP gateway 10 by call agent 15, and henceforth gateways 10 and 20 can communicate using the respective IP addresses by employing, for example, Real-Time Protocol (RTP).

Although the FIG. 1 arrangement depicts VoP gateways 10 and 20 coupling packet network 300 to respective Public Switched Telephone Networks (PSTNs) 100 and 200, they can be connected directly to user devices such as telephones. The functionality of a VoP gateway can even be embedded in devices to form packet phones or integrated packet-circuit voice integrated switching systems. When embedded in Customer Premises Equipment such gateways are sometimes called Media Terminal Adapters (MTAs). These can also be called untrusted end points. Call agents are sometimes called Call Servers or Call Proxy Servers.

When there are multiple call agents in a network arrangement, as shown in FIG. 1, each one typically communicates with a subset of gateways under its control. Each of these subsets is a *domain*. When it is desired to set up a call between domains, for example, domains 306 and 307, the respective call agents communicate with each other, as described above.

In the above example, network 300 was chosen to employ the Internet Protocol (IP), but it should be understood that Asynchronous Transfer Mode (ATM), Frame Relay (FR) or any other packet protocol that is suitable for transporting voice packets may be employed. The call set-up procedure for non-IP packet networks is similar to the procedure outlined above for IP networks.

A highly desirable characteristic of the FIG. 1 arrangement is the separation of *Call Control* from *Connection Control*. In this model, the techniques, signaling messages, procedures, etc., used to establish the logical voice connection between end-users is independent of the techniques, signaling messages, procedures etc., used to establish the connection that carries the voice packets in the packet network. In this way, customers can have and retain the same voice features regardless of whether the underlying transport technology is circuit-



as any special call features may be centralized. In addition, voice is usually billed on a duration basis, not a packet basis, and the packet network has no knowledge of call duration. Therefore, it is required that gateways 10 and 20 (or corresponding MTAs) be allowed to send packets to each only when allowed by the call agents.

If, instead, one were to decide to bill on a packet usage basis, governed by the IP network, the gateways might use the call agent to exchange IP addresses but never use the IP network to exchange voice packets, preferring to use some other (cheaper) network. Therefore, even in the case of billing on a packet usage basis, it is required that there be an affirmative control by the call agent of the connections through network 300.

Another consideration is that, for security reasons, users may not want their "true" IP address to be disclosed to others. This is particularly true if a user is in a private network behind a proxy firewall.

One solution to this problem is presented in FIG. 2, where call agent 15 communicates with a special router 313 at the edge of domain 306 (via line 308), and call agent 25 communicates with special router 323 at the edge of domain 307 (via line 309). These special *edge switches* route packets only if they carry an IP address that was explicitly authorized by a call agent. In specifying the authorized IP addresses, the call agent is also able to specify the QoS level being paid for, and that provides the edge switches with information necessary to choose between packets that are to be routed vs. packets that are to be buffered, when the transmission load calls for buffering of some packets. To prohibit the gateways from being used without the packet network, the IP addresses are never communicated end to end. Call agent 25 maps the IP address that leads to user 21 into an arbitrary IP address and communicates the arbitrary/true IP address mapping to its edge switches. It then communicates the arbitrarily selected IP address to call agent 15 and, thence, to gateway 10. Similarly, call agent 15 maps the IP address that leads to user 11 into an arbitrary IP address and communicates the arbitrary/true IP address mapping of to its edge switches. It then communicates the arbitrarily selected IP address to call agent 25 and,

thence, to gateway 20. In this way, gateways 10 and 20 never know the true IP addresses of each other.

There are a number of problems with this solution.

- This solution requires precise timing between the packet network and the call agents. If the messages to the edge switches are sent too soon, customers can obtain free service (for a short duration); if too late, the voice path might not be established by the time gateway 20 is answered, resulting in clipping of the initial speech.
- The call agent must know the characteristics of the packet network, because the procedures for establishing connections are different for each type, and the packet network may provide permanent connections (PVCs), temporary connections (SVCs), or no connections at all (as in IP).
- An end-to-end connection may require several networks: private networks, local public networks, inter-exchange carrier networks, and/or international networks. This communication must take place in each of these separate networks, adding to the complexity.
- For reliability, it is desirable to have the option to serve a particular gateway by any one of a multiple number of call agents and edge switches. However, for any given call, only one specific call agent/edge switch pair is involved. Reliably establishing the communication between the right ones in real time is difficult and requires the call agents to have accurate knowledge of the connection network topology as well as either additional network elements to keep the status of each call agent and edge switch and/or some kind of broadcast mechanism to insure the "right" edge switch gets the information. Additionally, in some cases (e.g. failure), the connection may even be re-established in the middle of a call, again, preferably without interaction with, or even knowledge of, the call agent. The issue of reliability is further complicated by the distributed nature of most edge switches themselves, with termination cards within the edge switch performing much if not all of the connection processing. The connection request will be received by one termination card, necessitating the same communication needs as between the call agent/edge switch, in that either the correct card must be identified and

informed, or all requests must be broadcast to all cards.

### **Summary of the Invention**

The prior art problems are overcome and an advance in the art is achieved by eliminating the need for a call agent to send mapping information directly to edge switches. This is achieved by all edge nodes mapping received packet addresses in accordance with a predetermined function. The mapping according to the function may change at regular intervals, or upon the occurrence of specified events, and with each change, the communicating user is provided with a different address to be used. In one embodiment, the mapped destination address that is created is developed through a process that encrypts the true address. The changed mapping in the context of an encryption scheme can be effected by merely specifying a different random seed value in the encryption algorithm.

### **Brief Description of the Drawings**

FIG. 1 describes the prior art arrangement of establishing voice connections over a packet network;

FIG. 2 describes the use of back-to-back edge switches between networks that aim to insure no unauthorized transmissions between networks;

FIG. 3 shows an arrangement where edge switches perform mappings without direct communication from call agents;

FIG. 4 presents a signal flow diagram in conformance with the principles disclosed herein; and

FIG. 5 shows an arrangement with two networks interposed between the networks of the two communicating devices.

### **Detailed Description**

FIG. 3 illustrates a packet network arrangement that comports with the principles of this invention; and with respect to those principles, it is similar to the FIG. 1 arrangement. For illustrative purposes, however, instead of a single



network as shown in FIG. 1, FIG. 3 depicts an ATM network 310, and an ATM network 320; instead of gateway 10, PSTN 100 and user 11, there is an MTA 13 that is connected to switch 314 within network 310; and instead of gateway 20, PSTN 200 and user 12, there is an MTA 23 that is connected to switch 324 within network 320.

It is noted that the switches in ATM networks perform essentially the same function as do the routers in IP networks. In this disclosure, therefore, the term "node" is used to subsume both a router and a switch.

For convenience, it may be assumed that MTA 13 is in New York and network 310 is owned by Verizon, that MTA 23 is in Los Angeles and network 320 is owned by PacTel, and that the connection between networks 310 and 320 is either direct, or circuit-switched, over a network owned by an inter-exchange carrier (not shown). Also, MTA 13 homes-in onto edge switches 311 and 312, to illustrate that, for increased reliability, two parallel paths may be conditioned to carry a connection between MTAs 13 and 23. Likewise, MTA 23 homes-in onto edge switches 321 and 322.

In accord with the principles disclosed herein, edge switches of a packet network – being the only points of general entry from another packet network – translate a predetermined portion of the address of incoming packets in accordance with a predetermined functional mapping. The portion that is functionally mapped is that portion that is expected to have been previously mapped by another functional mapping. The portion that is not mapped is that portion that is considered to be "clear."

No information needs to be communicated from a call agent to its associated edge switches. This mapping may be employed in the edge switches of the entire network (e.g. network 310), in edge switches of a particular domain, in a particular edge switch of the network, or even associated only with a particular call. The mapping may be through operation of a specified functional expression, or table-based. Illustratively, the mapping may be a decryption of a value that, when decrypted, yields the address of the destination MTA.

Further in accord with the principles disclosed herein and in cognizance of

the actions taken at edge switches, a remote source of packets that arrive at the edge switches of a network and are destined to an MTA at a given network address of the network (or, expressed more generally, destined to a port that has a network address) is not provided with this given network address of the destination MTA but, rather, is provided with a *mapped version* of the given address. The mapped version of the given address is such that when processed by the edge switches (i.e., mapped/decrypted) results in the true network address of the destination MTA. For example, if the destination MTA has a network address  $j$ , the address provided to a remote source of packets is  $A \bullet Q(j)$ , where "A" corresponds to a concatenated address portion that is in the "clear," while the  $Q(j)$  is the mapped network address of the MTA. The remote source sends out packets that carry the address  $A \bullet Q(j)$ . Based the clear portion of the address, the packets reach the network where the desired MTA is located, and the edge switches apply the mapped portion of the incoming address,  $Q(j)$ , to function  $\mathcal{P}$ , to yield  $\mathcal{P}(Q(j))$ , which equals  $j$  because the functions  $\mathcal{P}$  and  $Q$  are chosen to have this property.

Because the call agent already knows the addresses of the MTAs in its domain, it is advantageous keep the mapping function  $Q(j)$  in the call agents.

The mapping that is carried out by the edge switches for general packet communication may be long-lived, or short-lived; for example, valid only for one minute. In applications where the mapping function  $\mathcal{P}$  is not fixed, the mapping function  $Q$  must change in synchronism with changes in mapping function  $\mathcal{P}$  (or vice versa). In applications where the changes occur based on time of day, for example, this can be achieved by use of a common clock. Illustratively, the changes in functions  $\mathcal{P}$  and  $Q$  might take place in response to a reception of a broadcast signal.

To illustrate further, a network might use a pair of complementary encryption keys for the functions  $\mathcal{P}$  and  $Q$  (i.e.,  $\mathcal{P}(Q(j))=j$ ). In such an arrangement, the remote MTAs are given an address that has been encrypted with the key that corresponds to  $Q$ , and the edge switches decrypt with the key that corresponds to  $\mathcal{P}$ . Both keys may be algorithmically developed using a starting value (sometimes

called a “seed”). For example, the arrangement between the call agents and the edge switches might be that both entities work off a common set of seed values that are respectively pre-stored in a memory of the call agent and in a memory of the edge switches, and each minute of the day they independently create their respective keys by accessing the same (or complementary) seed values. Encryption functions such as the ones described above are well known in the art. See, for example, “Applied Cryptography,” by Bruce Schneier, John Wiley & Sons, 1996.

The synchronization between the call agent’s interval clock when mapping function  $\mathcal{Q}$  is changed, and the clock interval mappings when the edge switches change the mapping function  $\mathcal{P}$  need not be precise and, therefore, there is no need for the call agent to communicate directly with the associated edge switches to insure this synchronization. Even for a relatively short time interval such as one minute, a time offset between the call agent and the edge switches of a few seconds is not a problem as long as the edge switches are quicker to switch to a new mapping function than the associated call agent, but continue to remember the old mapping function. Time-adjacent mappings can be selected so that a mapping of an address that was mapped in accordance with the immediately previous mapping function yields an address that is recognized to be incorrect. In such an event, the previous mapping function is used to produce the correct mapping.

FIG. 4 presents a signal flow diagram for an implementation in accord with the principles of this invention for the FIG. 3 arrangement. For this illustration, it is assumed that networks 310 and 320 are ATM networks using Bearer Independent Call Control (BICC) protocol for call agent to call agent signaling, and establishing Switched Virtual Circuits (SVCs) for connection control.

When MTA 13 wishes to place a call, it sends a service request to call agent (CA) 15 (line 101 – e.g., Q.2931 protocol). In sending the service request, MTA 13 provides information about its own network address, and the identity of the called party (for example, MTA 23). In response to the latter, call agent 15 queries its database (line 102) for the address of a call agent that handles the





[illegible]

Note that call agent 15 and call agent 25 need not have any knowledge of how packets are routed by edge switch 311 and edge switch 321. If conditions change and an edge switch fails or becomes congested, the other edge switches can route around these problems without any action or knowledge on the part of the call agents, as long as these edge switches have knowledge of the appropriate mapping functions. In some cases, this rerouting can be accomplished during the call when the packet protocol allows this, e.g., in the IP protocol, or some implementations of the ATM protocol. This rerouting can be accomplished at call setup without the call agents' knowledge of the connection topology and which specific edge switches will be involved in the call.

As indicated above, the connection between networks 310 and 320 can be direct, or through one or more networks. FIG. 5 explicitly illustrates this condition; with network 330 interposed between networks 310 and 320. To simplify the drawing, only one edge node is shown to be involved in the connection involving

a. Regression coefficients		b. Regression coefficients	
Dependent variable	Independent variable	Dependent variable	Independent variable
a. $\ln(\text{GDP per capita})$	$\ln(\text{GDP per capita})$	b. $\ln(\text{GDP per capita})$	$\ln(\text{GDP per capita})$
	$\ln(\text{GDP per capita})$		$\ln(\text{GDP per capita})$
	$\ln(\text{GDP per capita})$		$\ln(\text{GDP per capita})$
	$\ln(\text{GDP per capita})$		$\ln(\text{GDP per capita})$
b. $\ln(\text{GDP per capita})$	$\ln(\text{GDP per capita})$	c. $\ln(\text{GDP per capita})$	$\ln(\text{GDP per capita})$
	$\ln(\text{GDP per capita})$		$\ln(\text{GDP per capita})$
	$\ln(\text{GDP per capita})$		$\ln(\text{GDP per capita})$
	$\ln(\text{GDP per capita})$		$\ln(\text{GDP per capita})$
c. $\ln(\text{GDP per capita})$	$\ln(\text{GDP per capita})$	d. $\ln(\text{GDP per capita})$	$\ln(\text{GDP per capita})$
	$\ln(\text{GDP per capita})$		$\ln(\text{GDP per capita})$
	$\ln(\text{GDP per capita})$		$\ln(\text{GDP per capita})$
	$\ln(\text{GDP per capita})$		$\ln(\text{GDP per capita})$

When the intermediate networks wish to block traffic except that which they get paid for, one approach that can be employed is the functional mappings-inverse mappings that are disclosed herein. In accordance with this approach, the traversal through any network is preceded by a mapping of an address portion at the incoming edge node and, therefore, the “connect” message that MTA 23 needs to send in the FIG. 5 arrangement is addressed to

and the acknowledgement message is addressed to

The values  $A_{310}$  and  $\mathcal{Q}_{15}^{\ell}(X1)$  are provided to call agent 25 by call agent 15. Call agent 15 obtains the values  $A_{340}$ ,  $A_{330}$ , and  $A_{320}$  from its database, forwards values  $A_{310}$ , and  $A_{330}$  to call agent 45, and instructs it to send  $\mathcal{Q}_{45}^{\ell}(A_{310})$  and  $\mathcal{Q}_{45}^{\ell}(A_{330})$  to call agent 25. Similarly, call agent 15 forwards values  $A_{340}$ , and  $A_{330}$  to call agent 35, and instructs it to send  $\mathcal{Q}_{35}^{\ell}(A_{340})$  and  $\mathcal{Q}_{35}^{\ell}(A_{320})$  to call agent 25. Call agent 25 then provided MTA 23 with the above values, including  $\mathcal{Q}_{25}^{\ell}(X2)$ , thus supplying all of the necessary information for setting up a connection.

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knowledge of the path, is for each call agent to determine the next network in the path, map the previous network's address, and concatenate its "clear" address to the resultant address.

As indicated above, the selection of FIG. 3 network as an ATM network was merely illustrative. It should be noted that the principles disclosed herein are applicable to other packet technologies, call control protocols and connection methods.

It should be also appreciated that though the mappings performed in the edge nodes, and the mappings performed in the call agents are functional, in the sense that given an address the mapped value can be computed, this computing to obtain the mapped value can be replaced with a look-up table. It should also be appreciated that various, arbitrarily selected, parameters can be included in the process that chooses the mapping functions  $\mathcal{P}$  and  $\mathcal{Q}$ . This is particularly so when the call agents and the edge nodes take their respective cues for changing functions  $\mathcal{P}$  and  $\mathcal{Q}$  from a received broadcast signal.



## Claims

1. A method for communicating packets from a packet source in a first network to a packet destination in a second network, where said packet destination has a network address  $X$ , comprising the steps of:
  - communicating, from an element in said second network to an element in said first network, an address  $Y$  that corresponds to address  $X$  mapped with function  $\mathcal{Q}$ ;
  - and
  - mapping, in a node in said second network, at least a sub-field of an address field contained in packets received from said first network with a function  $\mathcal{P}$ , where  $\mathcal{Q}$  and  $\mathcal{P}$  are functions such that  $\mathcal{P}(\mathcal{Q}(X))=X$ .
2. The method of claim 1 where functions  $\mathcal{P}$  and  $\mathcal{Q}$  change upon occurrence of an event.
3. The method of claim 2 where said event is reception of a change-specification signal, or a specified change in the time-of-day.
4. The method of claim 2 where said functions  $\mathcal{Q}$  and  $\mathcal{P}$  change at regular time intervals.
5. The method of claim 1 where said changes to said mapping function  $\mathcal{Q}$  and mapping function  $\mathcal{P}$  are algorithmically determined.
6. The method of claim 1 where said changes to said mapping function  $\mathcal{Q}$  and mapping function  $\mathcal{P}$  are determined by reference to a table that is stored in said element of said second network, and a table that is stored in said node.
7. The method of claim 6 where said table in said node contains seed values that are used to develop a decryption function to serve as mapping function  $\mathcal{P}$ , and said table in said element of said second network contains seed values that are used to develop a decryption function to serve as mapping function  $\mathcal{Q}$ .



[illegible]

if said step of determining concludes that said result of said mapping does not correspond to a valid packet destination in said second network, mapping said at least a sub-field of an address field contained in packets received from said first network with a function  $\mathcal{P}'$ , which corresponds to the mapping function employed by said node prior to the last change in mapping function  $\mathcal{P}$ .

**20.** The method of claim 1 where said node in said second network receives said packets from said first network via one or more other networks.

mapping, in a node in said second network, at least a sub-field of an address field contained in packets received from said first network with a function  $\mathcal{P}$  that is a decryption function corresponding to said encryption function  $\mathcal{P}$ , and functions  $\mathcal{P}$  and  $\mathcal{Q}$  are generated independently of communication between said element in said second network and said node in said second network.

## **Abstract**

Control over the movement of packets is exercised by edge nodes of a network mapping the addresses of incoming packets in accordance with a prespecified functional mapping  $P$ . Remote sources of packets are provided address information that is mapped with a prespecified functional mapping  $Q$ , where mappings  $P$  and  $Q$  are such that  $P(Q(j))=j$ . The mapping change at regular intervals, or upon the occurrence of specified events, and with each change, the communicating the remote source is provided with a different mapped address to be used.

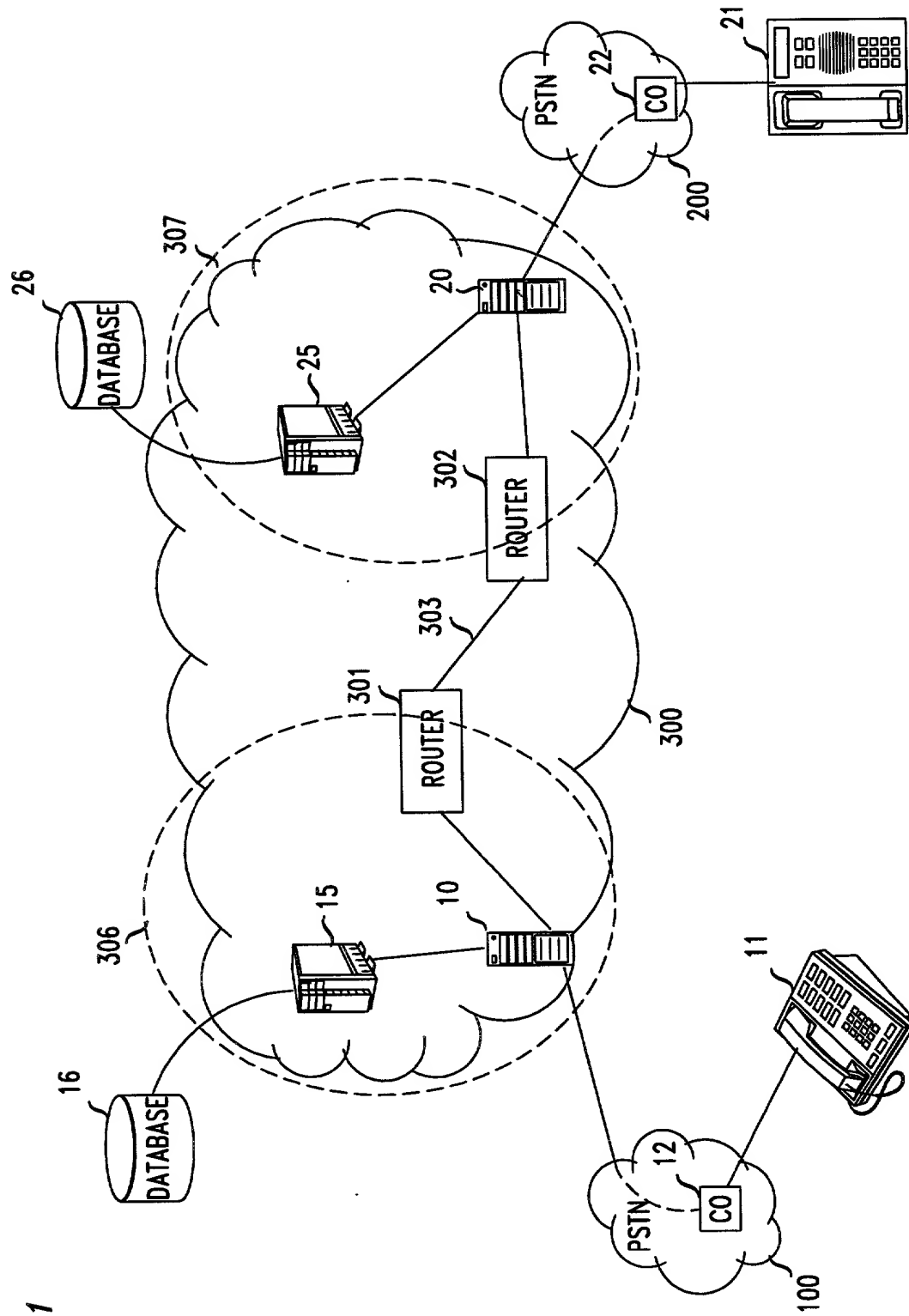


FIG. 1

000000 26549960

FIG. 2

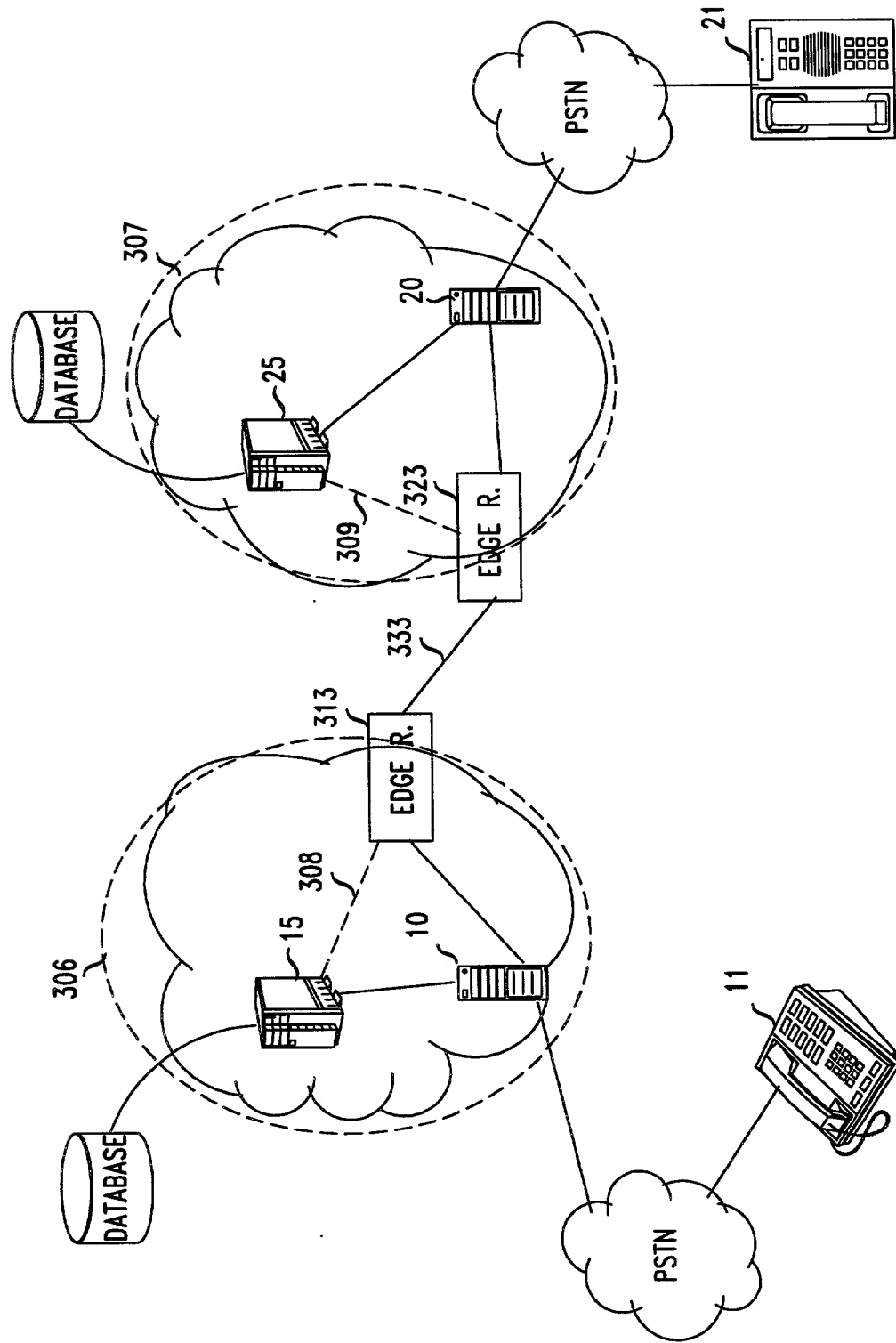


FIG. 3

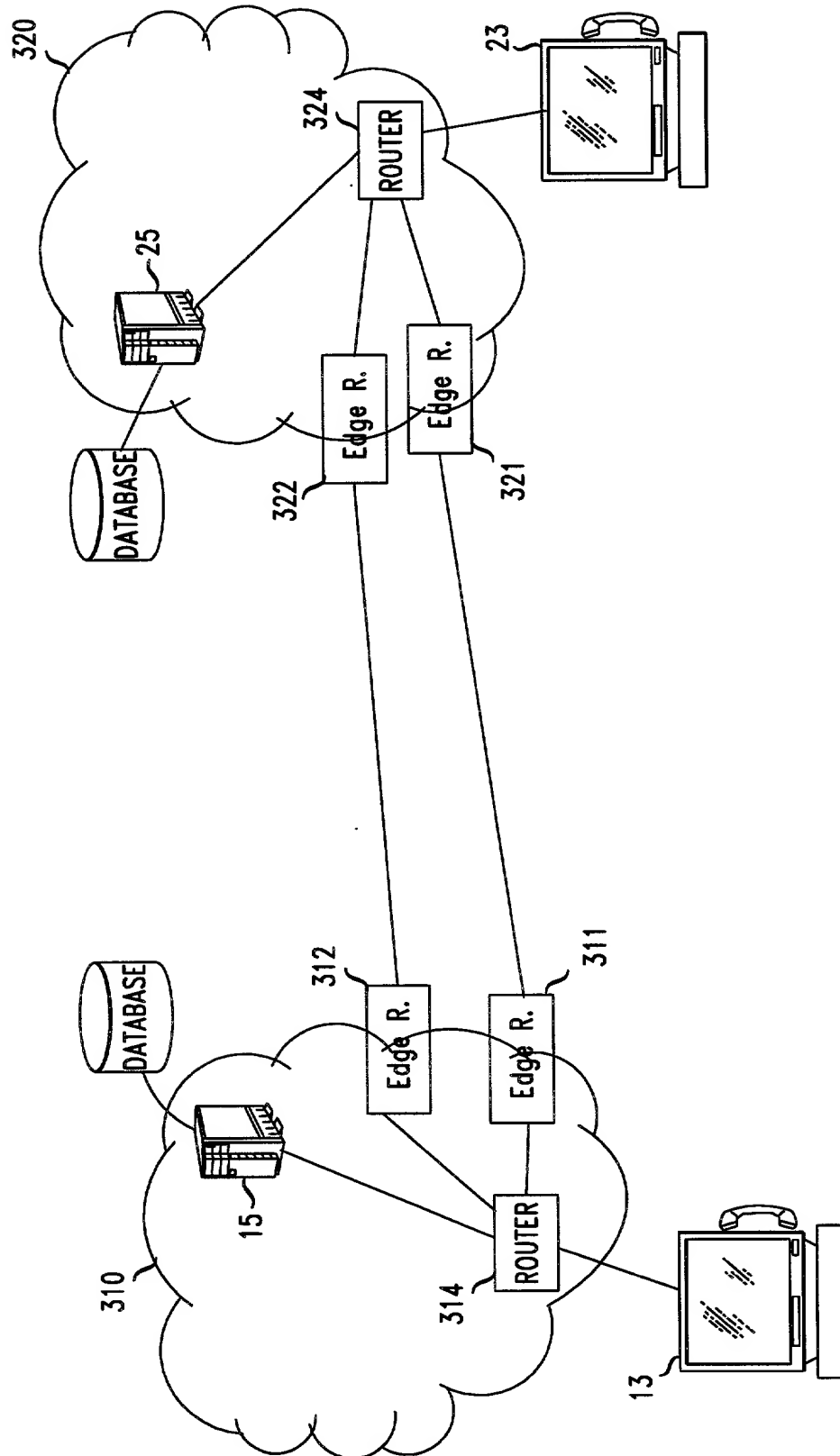


FIG. 4

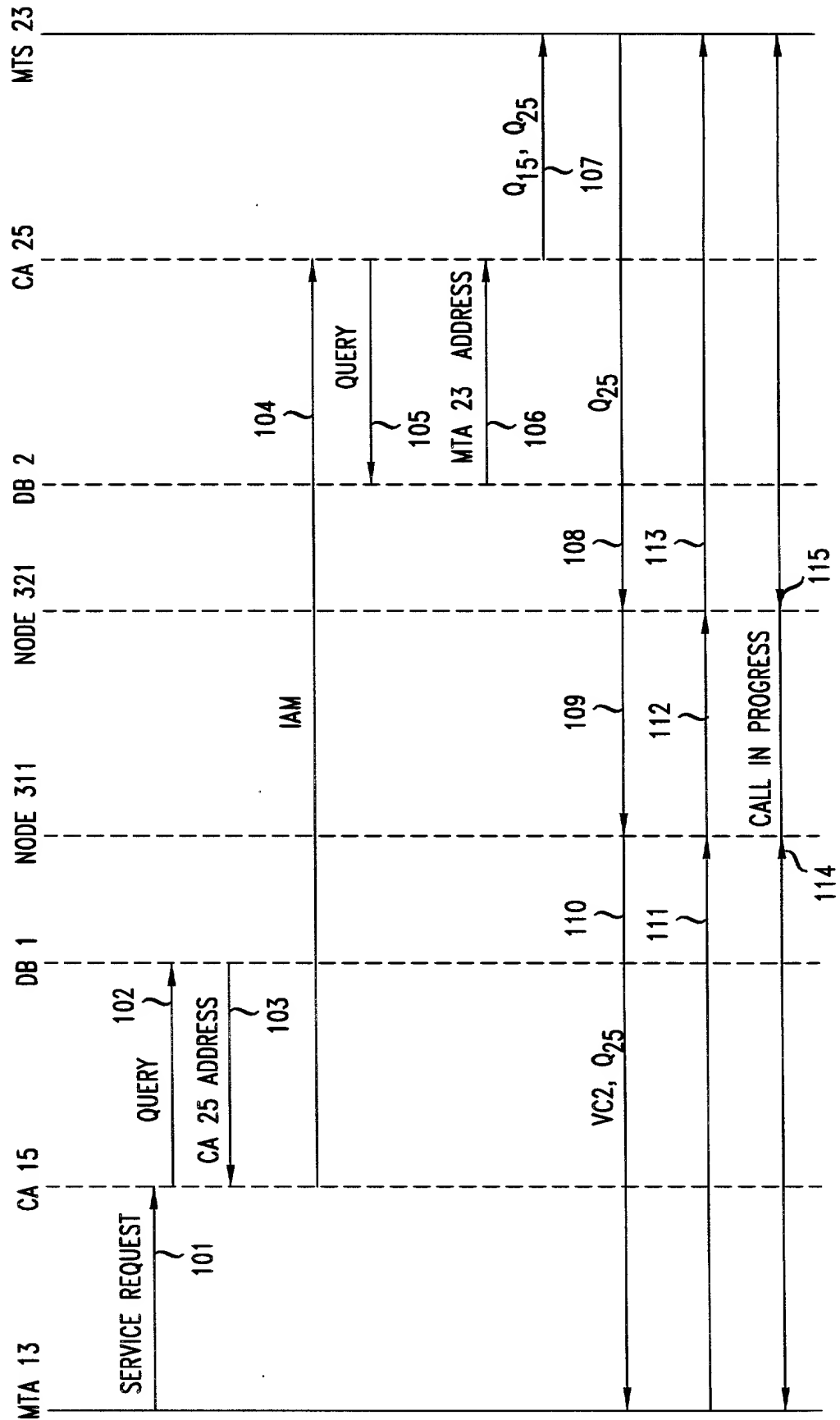
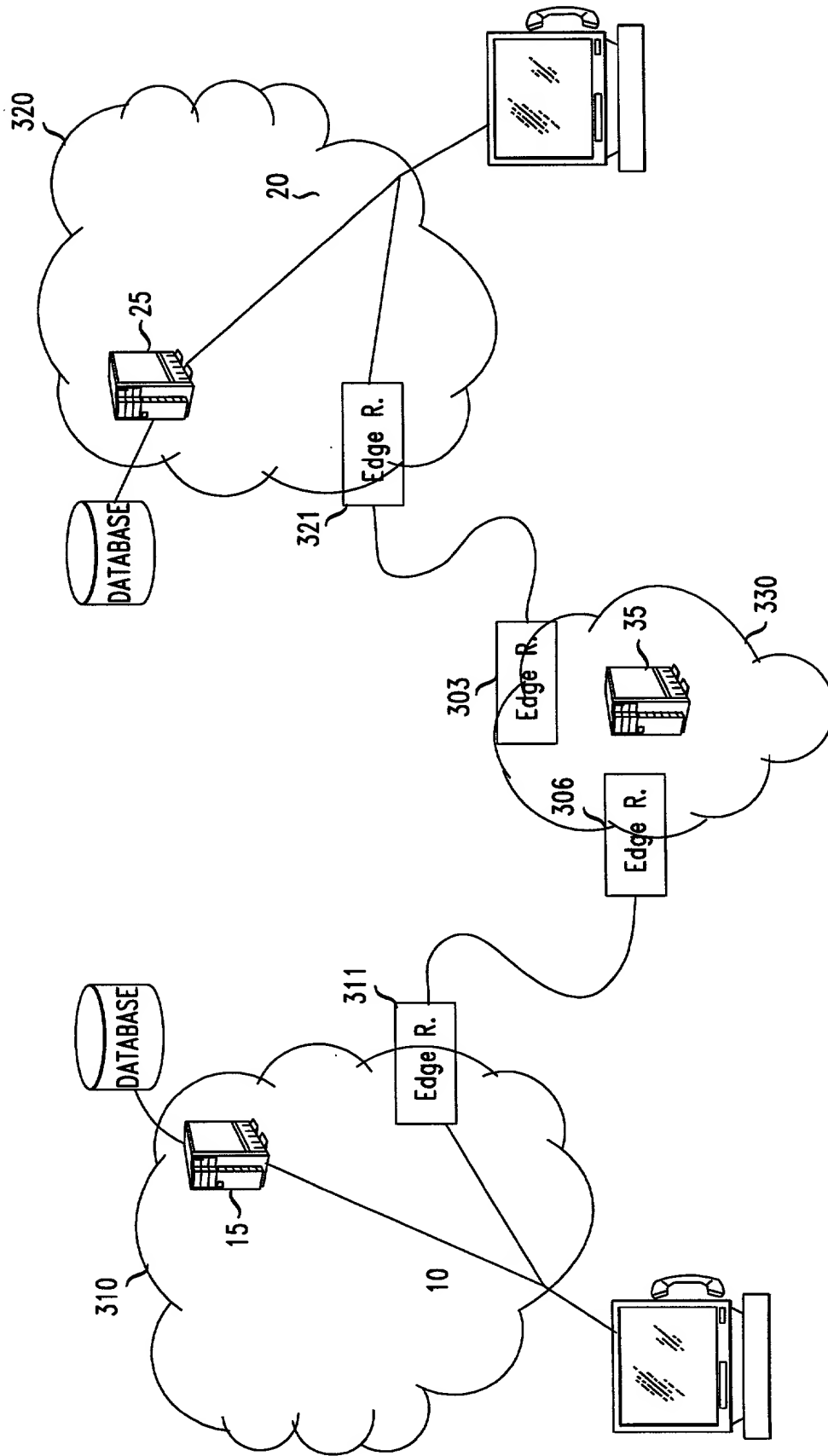




FIG. 5



## Declaration and Power of Attorney

My residence, post office address and citizenship are as stated below next to my name.

I believe I am an original, first and joint inventor of the subject matter which is claimed and for which a patent is sought on the invention entitled **Controlled Transmissions Across Packet Networks** the specification of which is attached hereto.

I hereby state that I have reviewed and understand the contents of the above-identified specification, including the claims, as amended by an amendment, if any, specifically referred to in this oath or declaration.

I acknowledge the duty to disclose all information known to me, which is material to patentability as defined in Title 37, Code of Federal Regulations, 1.56.

I hereby claim foreign priority benefits under Title 35, United States Code, 119 of any foreign application(s) for patent or inventors' certificate listed below and have also identified below any foreign application for patent or inventors' certificate having a filing date before that of the application on which priority is claimed:

None

I hereby claim the benefit under Title 35, United States Code, 120 of any United States application(s) listed below and, insofar as the subject matter of each of the claims of this application is not disclosed in the prior United States application in the manner provided by the first paragraph of Title 35, United States Code, 112, we acknowledge the duty to disclose all information known to us to be material to patentability as defined in Title 37, Code of Federal Regulations, 1.56 which became available between the filing date of the prior application and the national or PCT international filing date of this application:

None

I hereby declare that all statements made herein of my own knowledge are true and that all statements made on information and belief are believed to be true; and further that these statements were made with the knowledge that willful false statements and the like so made are punishable by fine or imprisonment, or both, under Section 1001 of Title 18 of the United States Code and that such willful false statements may jeopardize the validity of the application or any patent issued thereon.

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